

# GEMS OF TCS

## UNDECIDABILITY

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October 4, 2023

# ALAN TURING



1912–1954

# EVERYTHING IS A BIT STRING

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- Given input, algorithm
  - either eventually outputs some value
  - or never halts



# Halting Problem

# INFINITE LOOPS

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```
i = 0
while i <= 5:
    print('Infinite loop')
```

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```
i = 0
while i <= 5:
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```

---

```
x = True
while x:
    print('Infinite loop')
```

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  - The second input is string  $x$
  - $\text{HALT}(A, x) = 1$  if  $A$  halts on input  $x$
  - $\text{HALT}(A, x) = 0$  if  $A$  enters infinite loop on input  $x$



## APPLICATIONS OF HALTING PROBLEM

- Algorithm for HALT will help to design bug-free soft (and hardware)

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  - ...

Clearly, every function can be  
computed given sufficient time



Except this is **not** true

# HALTING IS UNDECIDABLE

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- Easy to solve for one input and one algorithm
- But impossible to solve for all inputs and algorithms
- Result holds for all computational models
- All non-trivial properties of algorithms are undecidable

Compiler

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# UNDECIDABLE PROBLEM

- Function  $A_{\text{diag}}(x)$  is defined as follows
- If the algorithm  $x$  on input  $x$  outputs 1, then  $A_{\text{diag}}(x) = 0$
- If the algorithm  $x$  on input  $x$  outputs other value or never halts, then  $A_{\text{diag}}(x) = 1$

# DIAGONALIZATION

PROOF

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## REDUCTION FROM DIAG TO HALT

- Assume there exists an algorithm for HALT
- Given input  $x$ , we check if the algorithm  $x$  halts on  $x$
- If it doesn't halt, output 1
- If it halts and outputs 1, output 0
- If it halts and outputs something else, output 1