

# MATRIX RIGIDITY

## RIGIDITY OF A FAMILY OF CIRCULANT MATRICES

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# KNOWN RIGIDITY BOUNDS

For CLB,  $R(\varepsilon_n) \geq n^{1+\delta}$

- $\mathcal{R}(r) \geq \frac{n^2}{r} \log \frac{n}{r}$

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- For linear  $r = \varepsilon n$ ,  $\mathcal{R}(\varepsilon n) \geq \underbrace{\Omega(n)}$

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- $\mathcal{R}(r) \geq \frac{n^2}{r} \log \frac{n}{r}$
- For linear  $r = \varepsilon n$ ,  $\mathcal{R}(\varepsilon n) \geq \underline{\Omega(n)}$
- Most limitations say that we can't achieve  $\mathcal{R}(\varepsilon n) \geq \boxed{n^{1+\delta}}$ .

$\forall \delta \exists \varepsilon \text{ s.t.}$

$$\mathcal{R}_n(\underline{n^{1-\delta}}) < n^{1+\delta}$$

$$n^{1+o(1)}$$

$$\underline{n \log \log n}$$

$$\underline{n \log^2 n}$$

$$\underline{n \cdot 2^{\sqrt{\log n}}}$$

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 $\mathcal{R}(\varepsilon n) \geq \underbrace{n^{1+\delta}}$   
*not acc  $n^{1+\sqrt{\delta}}$*
- What about just  $\mathcal{R}(\varepsilon n) \geq \underline{\omega}(n)$ ?
- Would suffice for circuit lower bounds against series-parallel circuits

# FAMILY OF CIRCULANT MATRICES

[CPR98] conjectured that an explicit family of circulant matrices has rigidity  $\mathcal{R}(\varepsilon n) \geq n(\log n)^\delta$ .

Last week:

No circulant matrix can

have  $\mathcal{R}(\varepsilon n) > n^{\delta + \epsilon}$

$$\left( \begin{array}{cccccc} a_1 & a_2 & a_3 & a_4 & \dots & a_n \\ a_2 & a_3 & a_4 & a_5 & \dots & a_1 \\ a_3 & a_4 & a_5 & a_6 & \dots & a_2 \\ a_4 & a_5 & a_6 & a_7 & \dots & a_3 \\ \vdots & \vdots & \vdots & \vdots & \ddots & \vdots \\ a_n & a_1 & a_2 & a_3 & \dots & a_{n-1} \end{array} \right) \leftarrow$$

[CPR98] describe  $a_1, \dots, a_n \in \{0, 1\}^n$

Only  $O(\log n)$  of  $a_1, \dots, a_n$  are ones, all the remaining entries are zeros.

In particular

$$\|M\|_0 = O(n \log n) \\ \Rightarrow R_M(r) \leq O(n \log n) \\ \forall r$$

Then conj  $R_M(n) \geq n^{(\log n)}$

# ODD ALTERNATING CYCLE CONJECTURE

## Conjecture

[CPR98] For every field  $\mathbb{F}$ , there exists an odd  $k$  and  $\varepsilon > 0$  such that every matrix  $M \in \mathbb{F}^{n \times n}$  with non-zeros on the main diagonal and  $\text{rk}(M) \leq \varepsilon n$  contains an cycle of length  $k$ .

# ODD ALTERNATING CYCLE CONJECTURE

Conjecture .

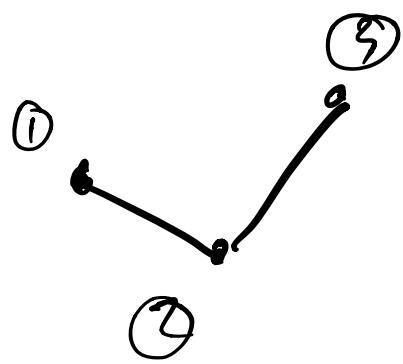
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Rigidity

$\overline{\text{over}} \mathbb{F}$

$\overline{\text{over}} \mathbb{F}$

[CPR98] This conjecture implies rigidity of an explicit circulant matrix



	1	2	3
1	1	1	1
2	1	1	1
3	1	1	1

	1	2	3	4	5
1	1	1	1	1	1
2	1	1	1	1	1
3	1	1	1	1	1
4	1	1	1	1	1
5	1	1	1	1	1

If  $\text{rk}(M) \leq \epsilon n$ , then

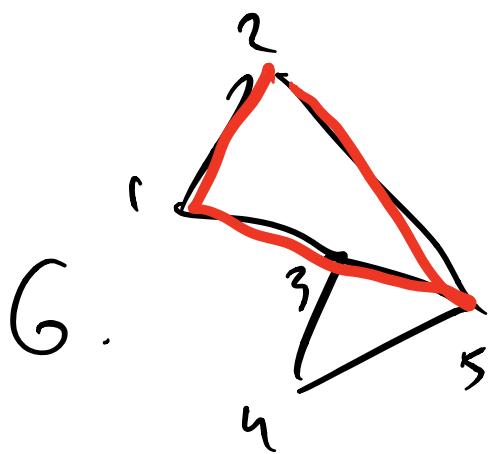
$G$  must contain a cycle  
of fixed odd length.

$G$  without short cycles (of odd length)  
must have high rank

IF  $\text{rk}(M) \leq \frac{n}{6} \Rightarrow$   
G must contain a 4-cycle

For every even length of a cycle

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	1	2	3	4	5
1	1	0			
2		1	1		
3			1	1	1
4				1	1
5					1

Conj. 1 IF  $\text{rh}(M) \leq \epsilon n$

$\Rightarrow G$  must contain a triangle

[CPR98]: Counter example:

Found graphs w/  $n^{2/3}$   
and without triangles.

[CPR98]. New Conjecture:

$\exists$  odd  $k$ ,  $\epsilon > 0$  s.t.

IF  $\text{rh}(M) \leq \epsilon n$

$\Rightarrow G$  must contain a cycle of length  $k$

# COUNTEREXAMPLE TO CONJECTURE

## Theorem (GH20)

For every odd integer  $k \geq 3$  there exists  $\delta > 0$  such that for every sufficiently large integer  $n$ , there exists a matrix  $M \in \mathbb{F}^{n \times n}$  with no odd cycle of length at most  $k$  such that for every finite field  $\mathbb{F}$ ,  $\text{rk}_{\mathbb{F}}(M) \leq n^{1-\delta}$ .

If  $k$  is odd we construct a graph with

(1)  $\text{rk}(M) \leq \underline{n^{1-\delta}}$

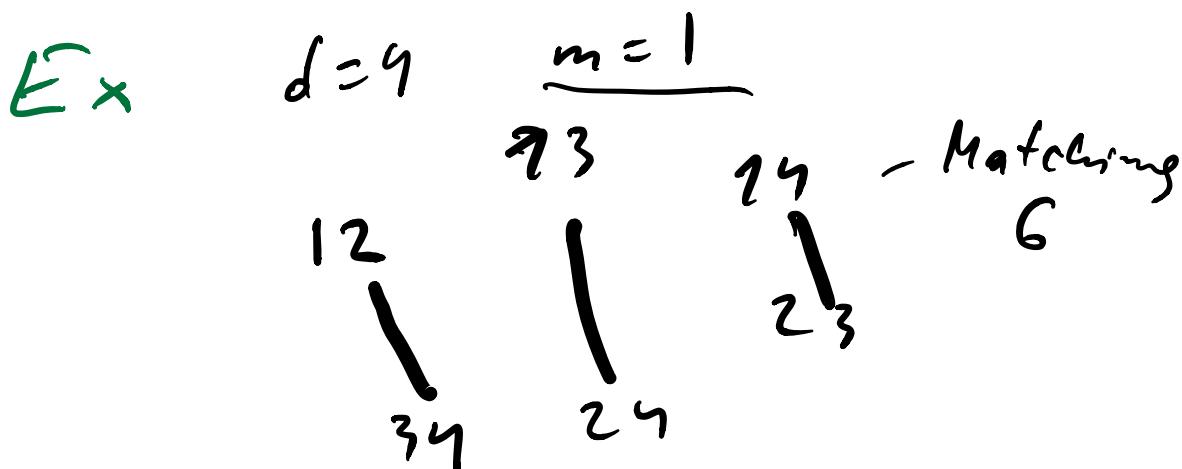
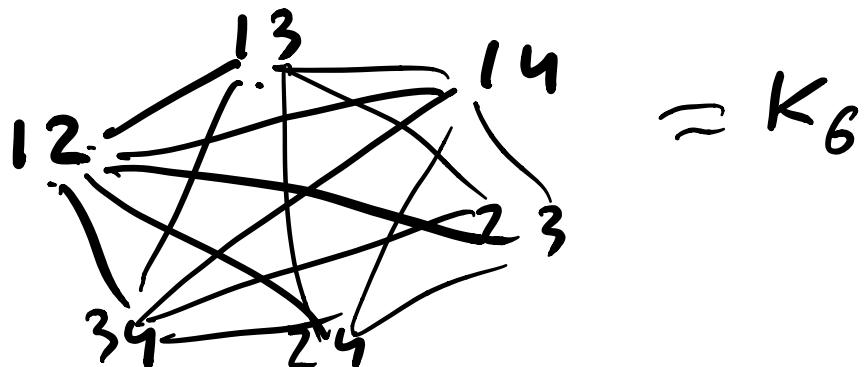
(2) no cycles of length  $k$  ...

Def Generalized Kneser graphs  
 $K(d, \underline{m})$ .

Vertices are all subsets of  $\{1, \dots, d\}$   
 of size  $= \frac{d}{2}$ .

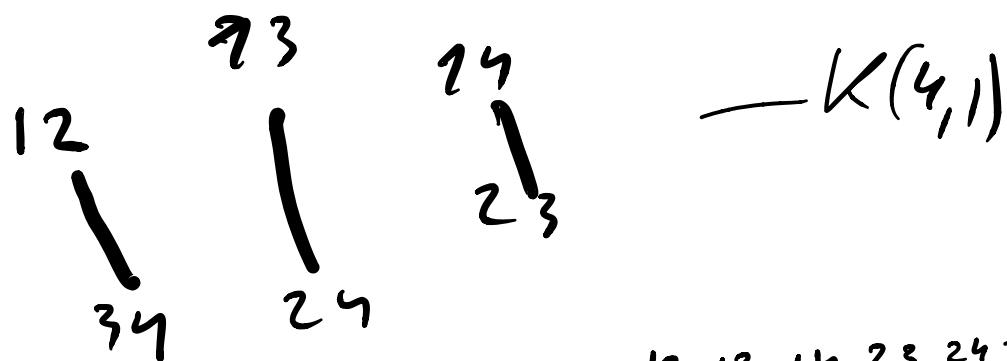
Two vertices  $I, J \subseteq \{1, \dots, d\}$   
 are connected iff  $|I \cap J| < m$ .

Ex  $d = 4 \quad m = 2$ .



Let's consider  $K(d, m)$

Let  $M$  be its adjacency matrix  
with 1s on the main diagonal.



$K(d, m)$

$$M = \begin{array}{c|cccccc} & 12 & 13 & 14 & 23 & 24 & 34 \\ \hline 12 & 1 & & & & & \\ 13 & & 1 & & & & \\ 14 & & & 1 & & & \\ 23 & & & & 1 & & \\ 24 & & & & & 1 & \\ 34 & & & & & & 1 \end{array}$$

$$\boxed{\text{rk}(M) \leq \sum_{i=0}^{\frac{d-m}{2}} \binom{d}{i}}$$

$$K(d, m) \\ \# \text{vertices} = \binom{d}{d/2} \approx \frac{2^d}{\sqrt{d}} = 2^{d - o(d)}$$

$$m = \epsilon^d \\ Rk(M) = \sum_{i=0}^{d(\frac{1}{2}-\epsilon)} \binom{d}{i} \leq \\ \leq 2^{d H(\frac{1}{2}-\epsilon)} \approx 2^{d(1-\epsilon^2)} \\ = \boxed{2^{d(1-\delta)}}$$

$$\# \text{vertices} \text{ is } N \approx 2^d$$

$$Rk \text{ is } \leq N^{(1-\delta)}$$

$$m = \varepsilon d < \frac{d}{2}$$

$$\begin{aligned} |I| = |J| &= d/2 \\ I, J &\subseteq \{1, \dots, d\} \end{aligned}$$

	12	13	14	23	24	34
12	1					
13		1				
14			1			
23				1		
24					1	
34						1

$$M_{I,J} =$$

$$= f(I, J) = \boxed{\sum_{i=m}^{d/2-1} (|I \cap J| - i)}$$

Proof: (1) Green ones - on the main diag.

$$M_{I,I} = f(I, I) =$$

$$= \sum_{i=m}^{d/2-1} (d/2 - i) \neq 0$$

$$(2) M_{I,J} = f(I, J) = 0$$

$$\text{IFF } |I \cap J| \geq m$$

$$= \boxed{\prod_{i=m}^{d/2-1} (|I \cap J| - i)}$$

$$\frac{m \leq |I \cap J| \leq \frac{d}{2}-1}{|I|=|J|=\frac{d}{2}} \Rightarrow |I \cap J| \leq \frac{d}{2}-1$$

$M_{I,J}=0$  IFF  $|I \cap J| \geq m$

$M_{I,J} \neq 0$  IFF  $|I \cap J| < m$


  
 dof of an edge in the  
 Kneser graph  $K(d, m)$

$P_{x,y} = P(x,y)$  where

$P$  has  $m$  monomials  $\Rightarrow$

$$\text{rk}(P) \leq m$$

Our  $\deg\left(\frac{d}{2}-1\right)-m+1 \approx \frac{d}{2}-m$

# monomials

$$\sum_{i=0}^{\frac{d}{2}-m} \binom{d}{i}$$

$$\Rightarrow \text{rank}(M) \leq \sum_{i=0}^{\frac{d}{2}-m} \binom{d}{i}$$



$$= 2^{dH\left(\frac{1}{2}-\epsilon\right)} = 2^{d(1-\delta)}$$

Sublinear in  $H$   
of vertices

$k$ -Fixed odd constant

$$m = \frac{d}{2^k}, \text{ consider}$$

$K(d, m)$

Claim  $\rightarrow$  has no cycles of  
odd length  $\leq k$ .

$$k=3$$

$K(d, \frac{d}{6})$

$$|I \cap J| < \frac{d}{6}$$

Assume  $I - J - L$

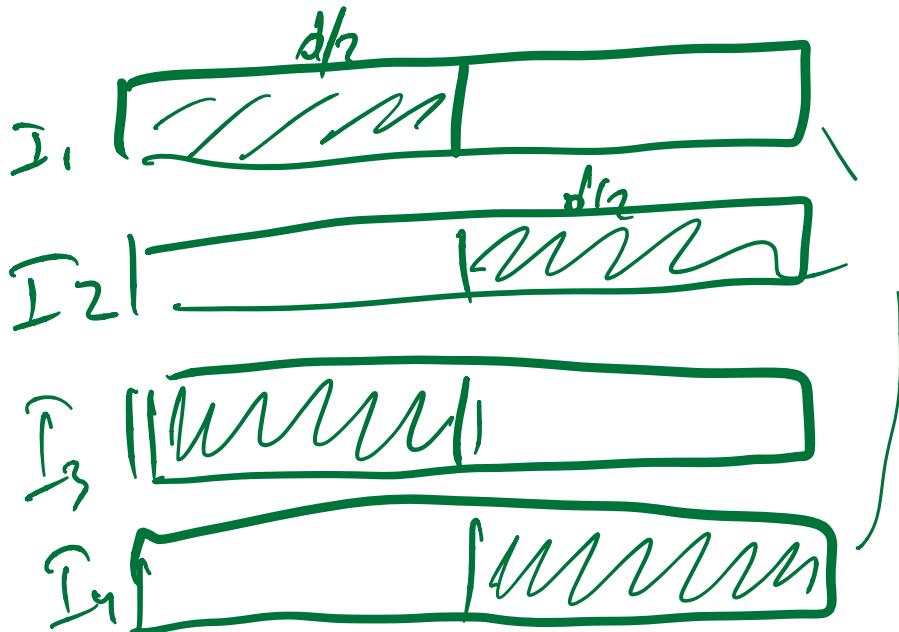
$$|I|=|J|=|L|=\frac{d}{2}$$

$$|I \cap J|, |I \cap L|, |J \cap L| < \frac{d}{6}$$



Even length

$$I_1 - I_2 - I_3 - I_4 - I_5 - I_6$$



If  $\kappa$  is any odd number  
then  
 $|I_1 \cap I_\kappa|$  is large — they are  
not connected  $\Rightarrow$  no cycles of  
length  $\kappa$ .